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Zhu, Qianqian; Panichella, Annibale; Zaidman, Andy

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An Investigation of Compression Techniques to Speed up Mutation Testing

Qianqian Zhu  
Delft University of Technology  
Email: qianqian.zhu@tudelft.nl

Annibale Panichella  
University of Luxembourg  
Email: annibale.panichella@uni.lu

Andy Zaidman  
Delft University of Technology  
Email: a.e.zaidman@tudelft.nl

Abstract—Mutation testing is widely considered as a high-end test coverage criterion due to the vast number of mutants it generates. Although many efforts have been made to reduce the computational cost of mutation testing, in practice, the scalability issue remains. In this paper, we explore whether we can use compression techniques to improve the efficiency of strong mutation based on weak mutation information. Our investigation is centred around six mutation compression strategies that we have devised. More specifically, we adopt overlapped grouping and Formal Concept Analysis (FCA) to cluster mutants and test cases based on the reachability (code coverage) and necessity (weak mutation) conditions. Moreover, we leverage mutation knowledge (mutation locations and mutation operator types) during compression. To evaluate our method, we conducted a study on 20 open source Java projects using manually written tests. We also compare our method with pure random sampling and weak mutation. The overall results show that mutant compression techniques are a better choice than random sampling and weak mutation in practice: they can effectively speed up strong mutation 6.3 to 94.3 times with an accuracy of >90%.

I. INTRODUCTION

Mutation testing has been actively investigated as a technique to evaluate the quality of test suites [1]. The main idea is (i) to introduce small syntactic changes (mutants) into the production code using mutation operators, and (ii) to measure the ability of a given test suite in detecting them [2]. One of the benefits reported in literature is that mutation testing provides a better measure of the fault detection capability of test suites compared to other test coverage criteria [3]–[5]. Despite its well-known advantages, mutation testing remains an extremely expensive activity since it requires to re-run the test suites against each mutant, whose number increases exponentially with the size of the program under test [6].

To address this limitation, several methods have been proposed and these can be classified in three main categories [7]: (do fewer) selecting fewer mutants to evaluate [8], [9], (do smarter) using run-time information to avoid unnecessary test executions [10], [11], (do faster) reducing the execution time for each single mutant [12]. Techniques falling into the first category are the most investigated. Indeed, researchers have proposed various strategies to sample mutants, such as random sampling [13], mutation operator selection [9], clustering [14], static analysis [15], and machine learning based sampling [16].

Recently, Gopinath et al. [17] have challenged the effectiveness and efficiency of mutation reduction strategies: their empirical evaluation with eight common mutant reduction techniques showed that none of them provide any practical advantage over pure random sampling. Although some techniques showed small improvements in effectiveness, the gains do not compensate for the extra overhead. Therefore, there is a need for reduction techniques that are not only more effective, but also more efficient compared to random sampling.

This paper originates from the insights of Gopinath et al. [17] and focuses on the mutant reduction technique recently proposed by Zhu et al. [18]. Zhu et al. tackle the problem of reducing the cost of mutation testing by combining do fewer and do smarter techniques through data compression methods. First, weak mutation (do smarter) is used to determine which mutants lead to an infection state through one single execution of the test suite against the original program. Then, formal concept analysis [19] (FCA) is applied to derive the maximal groupings [18], which are two-way clusters of mutants and tests. Each maximal grouping is composed of a set of mutants $M$ and a set of tests $T$ with the property that any mutant in $M$ is weakly killed by any test in $T$. Finally, a do fewer strategy is applied by running one single test case (test selection) against one single mutant (mutant sampling) from each maximal grouping [18]. Their empirical study with five Java programs and automatically generated unit test suites showed that FCA reduces the execution time of mutation testing by up to 85%.

In this paper, we spot two important limitations of FCA that can affect its ability to correctly estimate strong mutation. First, FCA groups mutants and tests according to weak mutation only: mutants leading to an infection state when running the same test case $t$ (or set of tests) are assumed to be redundant. This is why only one mutant in each maximal grouping is evaluated for strong mutation. However, mutants that are redundant in terms of weak mutation are not necessarily redundant in terms of strong mutation, because, for example, they are injected in different code locations (e.g., different methods) or are generated by different mutation operators. Second, Zhu et al. [18] focused only on maximal groupings, which may leave some tests and/or some mutants not assigned to any maximal grouping. As consequence, the estimated mutation score may be inaccurate.

To overcome these limitations, we enhance FCA with (i) mutant location and (ii) mutation operator type information when grouping mutants and tests according to weak mutation. This prevents mutants infecting different statements or generated by different operators to be inserted in the same grouping.
We also investigate maximal and non-maximal groupings to prevent final clusters from missing test cases and/or mutants.

To evaluate the benefits of our enhancements, we conducted an empirical study with 20 open-source Java projects and using the test suites manually written by the original developers. Then, we compare the different variants of the FCA-based technique with and without our enhancements and against weak mutation and pure random sampling (with 10% as sampling percentage).

Our results show that FCA with our enhancements is more accurate in estimating the (strong) mutation score compared to (i) the original FCA-based technique by Zhu et al. [18], (ii) random sampling, and (iii) weak mutation. In particular, we find that the compression strategy based on non-maximal groupings and enriched with mutant location information (referred as overlap+mloc in the remainder of the paper) estimates the strong mutation score with an average absolute error of 5% and an average accuracy of 93% while being five time faster than strong mutation, i.e., its average speed up is 5X. Instead, random sampling achieves a higher absolute error of 13% while requiring the same execution time of overlap+mloc, i.e., its average speed up is 5X as well. The other compression strategies lead to larger speed-up scores (up to 18X) but with the cost of having a larger absolute error, which ranges between 5% (i.e., the absolute error of overlap+mloc) and 13% (i.e., the absolute error of random sampling) on average. Therefore, our findings challenge prior results [17] as we find that mutation strategies based on compression methods (and FCA in particular) are more effective and/or more efficient than random sampling. Finally, weak mutation is the fastest technique but it also produces the largest absolute error of 23% on average.

As a final remark, we observe that all FCA-variants allow to estimate whether each individual mutant is strongly killed or not based on relatively few test executions by relying on the two-way clusters generated with FCA. Instead, random sampling does not take into account relationships between mutants and test cases, possibly also leading to underrepresented areas of production code in the estimation.

II. BACKGROUND AND RELATED WORK

In this section, we begin with an overview of mutation reduction techniques in literature and the works that motivate our approach. Then, we introduce the crucial concepts and theories on which our approach is based.

A. Mutation Reduction Strategies

Techniques for reducing the high computational cost have been an active area of research. Offutt and Untch [7]’s literature review summarises these approaches into three categories: do fewer, do smarter and do faster. The most well-known techniques for reducing the computational cost of mutation testing are random sampling [8], selective mutation [9], weak mutation [10] and mutant schema [20]. The aforementioned methods are independent of the program under test which can be flexibly combined with our methodology.

More recently, researchers have aimed to make further gains by including run-time information; a widely-adopted strategy is to execute the test suite on the original program before mutation execution to avoid unnecessary executions. Coverage-based optimisation filters out test executions when a test case does not cover the mutated statement; this optimisation is in use in existing tools such as JAVALANCHE [21], Major [22] and PIT/PITest [23]. Infection-based optimisation on the other hand only executes a test case on a mutant when the test infects the state of the mutant, filtering out weakly live mutants.

Just et al. [11] improved upon this by only executing a test on a mutant if the execution state of the mutated expression propagates to a top-level expression; they also partitioned mutants based on their intermediate results. Ma and Kim [24] applied a similar idea to cluster mutants for each test case by comparing the values of innermost expressions. Compared to Just et al. [11] and Ma and Kim [24]’s, we partition mutants for all test cases instead of targeting each test case.

Mutant clustering’s aim is to reduce the number of mutants based on the similarity of mutants instead of random sampling. Hussain [14] applied clustering algorithms (e.g. K-means), however, the approach requires the execution of all mutants against all the test cases, which cannot reduce the overhead during the mutation execution. Later, Ji et al. [13] measured the similarity of the mutants using domain analysis. They divide mutants based on static control flow analysis. But they only manually analysed the clustering accuracy without indicating the runtime overhead caused by the domain analysis. Different from these works, our approach groups mutants based on their reachability and necessity conditions against the tests.

An approach that eliminates redundant mutants is mutant subsumption (e.g. [25]–[27]). However, mutant subsumption requires full knowledge of the mutation kill matrix, which requires the execution of every mutant against every test. Computationally this process is more costly than traditional strong mutation, thus cannot be used to speed-up mutation execution. Furthermore, test prioritization and reduction are also used to speed up mutation testing, e.g. Zhang et al. [28].

Moreover, Zhang et al. have recently proposed Predictive Mutation Testing (PMT) to predict mutation testing results without execution [29]. They extracted 12 features from the programs and constructed a classification model to predict whether a mutant is killed or surviving. Their experiment showed that PMT could improve the efficiency of mutation testing by up to 151.4 times with a small loss in accuracy. Despite high efficiency, their approach needs to collect a series of program features, which requires different tools to fulfil; this is a substantial burden for the common programmer. Unlike Zhang et al. [29], we do not require any additional program features; the weak mutation information needed for our mutant clustering and data compression can be collected by our tool during the initial execution against the original program.

Our approach is an extension of our previous work [18]. Despite the encouraging results, we have identified two important limitations of our initial methods as mentioned before (in Section I), i.e., (1) weak mutation information is not enough;
and (2) FCA could lead to missing mutants and/or tests. To address these limitations, we propose another compressing strategy, i.e., overlapped grouping, which is the simplest and strictest clustering method. Moreover, we take full advantage of mutation location and operator type knowledge when compressing.

B. Mutant Compression

We now describe the core concepts behind our approach: weak mutation and FCA-based compression technique [18].

Weak Mutation. For a test case \( t \) to kill a mutant \( m \) which mutates the statement \( s \) of a program \( P \), there are three conditions [30]: (i) reachability: the execution of \( t \) must cover \( s \); (ii) necessity: the execution state of \( m \) is different from the execution state of \( s \); (iii) sufficiency: the incorrect state of \( m \) must propagate to the output causing a failure in \( t \).

Weak mutation uses the necessity condition, i.e., a mutant is killed if its execution leads to a state change. For example, the expression \( c=a+b \) and its mutated version \( c=a/0 \) have different outcomes (i.e., the mutant is weakly killed) if \( a \neq 0 \). Differently from strong mutation, weak mutation scores can be computed with one single execution of each test by instrumenting the mutated locations [31].

FCA-based compression technique [18]. Formal Concept Analysis (FCA) was originally a data analysis method and has shown to be a powerful mathematical technique to convey and summarize large amounts of information [19]. It takes as input the formal context which is a structure \( C = (O, A, I) \) where \( O \) is the set of objects, \( A \) is the set of attributes while \( I \subseteq O \times A \) is a binary relation between \( O \) and \( A \). Then, FCA produces the concept lattice, which is a collection of formal concepts in the data ordered by sub-concept relations, i.e., from super-concepts to sub-concepts. Each formal concept is composed of (i) a group of objects sharing the same attributes, and (ii) all attributes that apply to the objects in the concept [19].

In mutation testing context, the objects in \( O \) are the mutants, the attributes in \( A \) are the test cases, and \( I \) is the mutant-by-test infection matrix. Then, FCA derives formal concepts that represent groups of mutants that are weakly killed by the same subset of tests. In other words, the output of FCA can be viewed as two-way clustering since mutants and tests are grouped in concepts such that all mutants in the same concept are weakly killed by all tests in that concept. Among these concepts, FCA-based compression technique only considers the maximal concepts that are directly connected to the exit point in the lattice hierarchy that are referred to as maximal groupings.

After obtaining the maximal groupings from the concept lattice, this approach first compresses the mutant-by-test infection matrix by condensing the rows, i.e., select one mutant from each maximal grouping. Then to further perform the compression on the columns, there are three approaches for test case selection: (i) random: randomly select one test from each maximal grouping; (ii) Set cover based: find a sufficient subset of test cases that weakly kill all possible mutants. i.e., at each stage, choose the test that weakly kills the largest number of uncovered maximal groupings; (iii) Sorting-based: select the test cases with the largest number of maximal groupings at each stage until all possible mutants are covered.

However, as our previous study [18] show, applying test case selection on FCA groupings leads to a relatively small (5.89%) reduction in execution time. Therefore, in this paper, we do not use test case selection. Instead, we apply dynamic coverage-based optimisation [21–23] and infection-based optimisation [11], [24] (see Sec. II) to filter out unnecessary executions.

III. APPROACH

A. Overall Methodology

Our 6-step compression strategy is illustrated in Figure 1:

1. Instrumentation. We instrument the original program to keep track of the mutation locations: at every mutation point we insert all the mutants (mutated codes) right after the original one and assign a unique id to each mutant for later activation (we applied the technique of the mutant schemata [12]). To perform weak mutation, we also insert the comparison instructions at each mutation point to compare the intermediate states of the original program and mutated part (we compare the state after the first execution of the innermost expression that surrounds the mutant). We insert additional instructions to record information of each mutant including its location, operator type and mutation details (e.g., \( m_1 \) on Line 12 applies replace constant operator: \( 0 \rightarrow 1 \)).

2. Test execution. Once instrumented, the test suite is executed once on the original program. During this stage, we record the mutants that are touched by the tests, as well as the ones which are weakly killed by the tests. Only the instructions related to weak mutation and mutant information collection are executed at this stage. No mutants are activated.

3. Reachability and necessity analysis. The results of the previous stage are stored in the mutant-by-test reachability and mutant-by-test necessity matrices. Let \( P \) be the program under analysis and let \( T \) be the test suite; let \( M \) be the set of mutants for the program \( P \) generated by preselected mutation operators. A mutant-by-test reachability matrix is a \( m \times n \) matrix where \( m \) is the number of mutants, \( n \) is the number of test cases in \( T \), and an entry \( x_{i,j} \) is a binary value indicating whether the statement containing the \( i \)-th mutant is executed \( (x_{i,j} = 1) \) or not \( (x_{i,j} = 0) \) by the \( j \)-th test \( \in T \). The mutant-by-test necessity matrix has the same size as the mutant-by-test matrix.

Fig. 1. Overall Methodology of Mutation Compression Strategies
reachability matrix, but the binary entry \( x_{i,j} \) represents the outcome of weak mutation (\( x_{i,j} = 1 \) indicates weakly killed).

(4) Mutant clustering. Using the two aforementioned matrices, we apply clustering to group similar mutants together. We consider two clustering methods: (1) the overlapped grouping, and (2) FCA grouping from our previous study [18] (See Section II-B).

(5) Data compression. The mutant clusters are then used to compress the mutant-by-test matrix. The resulting matrix is likely to have lower dimensionality: the rows denote groups of mutants belonging to the same clusters; similarly, tests are grouped into clusters to form the columns. The compressed matrix is then used to apply mutants and execute tests for the strong mutation analysis. We take mutation knowledge into consideration during compression to achieve higher accuracy.

(6) Mutant Execution. The compressed matrix from the previous step is then used for the strong mutation analysis. Here, we load each mutant by its id and run the actual mutation execution against the tests.

The details of overlapped grouping and how we use mutation knowledge are described in the next sub-sections.

B. overlapped grouping

As mentioned earlier, the FCA-based compression technique could result in missing mutants and/or tests. Therefore, we propose the overlapped grouping to overcome this. The overlapped method is the simplest and strictest clustering method, i.e., elements are only grouped together if they are identical. Specifically, we first identify distinct mutants with regard to their reachability and necessity conditions against all the test cases. Subsequently, we group mutants having the same reachability and necessity conditions into one cluster.

The main difference between overlapped and FCA grouping is that overlapped grouping is stricter than FCA grouping. The overlapped grouping does not lose any information in the matrix, i.e., the clustering contains all the mutants. While FCA grouping loses some information since the main idea of FCA is to find the maximal sub-matrices (or formal concepts), and if a mutant or a test does not belong to a sub-matrix, FCA removes this mutant or test.

In our example in Figure 2, there are three maximal groupings, which are \{m5, m6\|t1, t2\}, \{m4\|t3\}, and \{m3\|t4\}. The other concepts in the lattice (e.g., \{m2, m5, m6\|t1\} in Figure 2) are already included in the maximal grouping by the sub-concept relation which is graphically represented by the hierarchy in the lattice. As for the overlapped grouping, it generates five clusters, i.e., \{m1\|t2\}, \{m2\|t1\}, \{m3\|t4\}, \{m4\|t3\} and \{m5, m6\|t1, t2\}. The example shows that for the clusters generated by the overlapped grouping no mutants are discarded, while the FCA grouping discards \( m1 \) and \( m2 \).

C. Mutation Knowledge

Once we have the mutant clusters, we compress the mutant-by-test infection matrix by condensing the rows, i.e., we select one/several mutants from each cluster to represent the whole group. To select the representative mutant(s), the FCA-based technique [18] adopted a random strategy which selects one mutant from each cluster at random, which we believe causes another limitation: based on weak mutation information alone, FCA could mis-cluster mutants at different code locations. Thus, we enhance FCA by adding mutation knowledge: (i) mutant location and (ii) mutation operator type.

More specifically, we investigate three mutant selection strategies: (1) random strategy; (2) random strategy with knowledge of the mutation operator type\(^1\); (3) random strategy with knowledge of the mutation location. The first one randomly chooses one mutant from each grouping/cluster as the representative mutant. The second strategy first divides each cluster into partitions by the type of mutation operator and then randomly selects one mutant from each partition. The third strategy partitions the cluster by the locations of the mutants (the line number) and then applies random selection; this guarantees that at least one mutant is selected for every potential mutation point. Notice that the second and third strategies might select more than one mutant from each grouping/cluster, which could lead to less speed-up in strong mutation.

To sum up, we devise one mutant clustering algorithm in addition to FCA as well as three mutant selection strategies, therefore, resulting in six compression strategies in total:

\[
\text{overlap} \quad \text{the combination of overlapped grouping and random strategy in mutant selection.}
\]

\[
\text{overlap+mop} \quad \text{the combination of overlapped grouping and random strategy with the knowledge of the mutation operator type in mutant selection.}
\]

\[
\text{overlap+mloc} \quad \text{the combination of overlapped grouping and random strategy with the knowledge of the mutation location in mutant selection.}
\]

\[
\text{fca} \quad \text{the combination of FCA grouping and random strategy in mutant selection.}
\]

\[
\text{fca+mop} \quad \text{the combination of FCA grouping and random strategy with the knowledge of the mutation operator type in mutant selection.}
\]

\[
\text{fca+mloc} \quad \text{the combination of FCA grouping and random strategy with the knowledge of the mutation location in mutant selection.}
\]

IV. EXPERIMENTAL STUDY

We conducted an empirical study to evaluate the effectiveness of the different compression strategies presented in the previous section. The goal of the study is to answer the following research questions:

- **RQ1:** How accurate are different compression techniques? We assess the ability of the six compression

\(^1\)Regarding the mutation operator type knowledge, we consider the operator type at a high level, e.g., arithmetic replacement operators.
strategies to estimate the strong mutation scores. We also assess their performance in comparison with random sampling and weak mutation.

- **RQ2**: How do compression techniques perform in terms of speed-up? We investigate the speed-up in terms of execution time that can be obtained when using each compression strategy over strong mutation. We also consider random sampling and weak mutation as baselines.

- **RQ3**: What is the trade-off between accuracy and speed-up for the compression techniques? We evaluate to what extent the compression strategies can reduce execution time while maintaining an accurate estimation of the strong mutation scores.

### A. Experimental Setup

To answer our research questions, we evaluated the six compression strategies using 20 open source projects publicly available on GitHub. Table I summarises the main characteristics of the selected projects. These projects have been randomly selected among the top 3000 GitHub repositories which (1) have most stars on 04/04/2017, (2) can be built using Maven, and (3) contain JUnit 4 test suites. In our study, we focus on the manually-written test suites available in the original project repositories.

As mentioned in Section III, we first need one test execution against the original program to collect statement coverage (i.e., the mutant-by-test reachability) and the weak mutation information (i.e., the mutant-by-test necessity matrix). To collect weak and strong mutation information, we implemented our own prototype tool. The instrumentation framework to generate mutants and detect the reachability and necessity condition is extracted from EvoSuite. Then, we integrated this instrumentation framework into our mutation testing runner. We record information about test cases (#id, method name, execution results, #touched mutants and #weakly killed mutants) and mutants (#id, mutation operator type, location and detailed information) for further analysis. After that, we run each test case against each mutant of the class under test (strong mutation) to establish the mutant-by-test sufficiency matrix which is used to evaluate our methods.

The mutation operators we adopted in this experiment are six method-level operators: replace arithmetic, replace bitwise, replace comparison, replace variable, replace constant, and insert unary. Further details about these mutation operators can be found in the paper by Fraser and Arcuri [31]. We opted for the mutation engine available in EvoSuite [31] because it instruments the production code at bytecode level and allows to directly measure the infection state for each mutant (weak mutation). To the best of our knowledge, no publicly-available mutation tool provides utilities for computing the weak mutation scores.

To answer the three RQs, we selected another two mutation reduction techniques (see Section II) for comparison: mutation sampling and weak mutation. We selected random sampling (do fewer strategy) as baseline because Gopinath et al. [17], [33] showed that none of the most common reduction strategies provide any practical advantage over pure random sampling. Moreover, we selected weak mutation (do smarter strategy) because it is one of the key components of all mutant compression techniques. Therefore, we considered it as an additional baseline to verify whether the other components of the compression strategies (e.g., computing the maximal groupings) are indeed needed. For random sampling, we set the sampling rate to 10% as suggested by Budd [6] and Acree [8]. They showed that 10% sampling could already estimate the mutation score with 99% of accuracy. It also corresponds to the sampling rate used by Gopinath et al. [17], [33]. Since random sampling and the mutant compression strategies involve random processing (i.e., in mutant selection), we carry out the corresponding random process 100 times for each project to address their randomised nature. In total, we compared eight mutation strategies: six compression strategies, random sampling (10%) and weak mutation.

### B. Evaluation Metrics

To answer RQ1, we selected two well-known performance metrics: the absolute error and the accuracy. Let \( M \) be a given mutation strategy (e.g., random sampling); let \( strong_M(C, T) \) be the percentage of mutants for a class \( C \) that are strongly killed by the test suite \( T \); let \( estimated_M(C, T) \) be the estimated percentage of mutants that are killed according to the strategy \( M \); the absolute error is defined as follows:

\[
AE(C, T) = \left| strong_M(C, T) - estimated_M(C, T) \right| \tag{1}
\]

The fifth column “COV” means the line coverage of the test suite which is measured by IntelliJ IDEA coverage runner.

The failures of the test cases in our tool is because the dependencies of these test cases need to be configured by Maven plugin; this cannot be solved in the current version of ComMT.

### Table I: Subject Programs

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<th>#Tests</th>
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<th>#Mutants</th>
<th>#Total Covered</th>
<th>#Weakly Killed</th>
<th>#Strongly Killed</th>
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<td>929</td>
</tr>
<tr>
<td>16</td>
<td>pel</td>
<td>3021</td>
<td>67</td>
<td>24</td>
<td>0.28</td>
<td>1176</td>
<td>380</td>
<td>294</td>
<td>205</td>
</tr>
<tr>
<td>17</td>
<td>stream-lib</td>
<td>4767</td>
<td>77</td>
<td>137</td>
<td>0.83</td>
<td>19007</td>
<td>9920</td>
<td>9445</td>
<td>6924</td>
</tr>
<tr>
<td>18</td>
<td>telegrambots</td>
<td>1480</td>
<td>21</td>
<td>30</td>
<td>0.20</td>
<td>772</td>
<td>221</td>
<td>196</td>
<td>52</td>
</tr>
<tr>
<td>19</td>
<td>vraport</td>
<td>12021</td>
<td>407</td>
<td>365</td>
<td>0.83</td>
<td>3490</td>
<td>2057</td>
<td>1652</td>
<td>7072</td>
</tr>
<tr>
<td>20</td>
<td>zt-zip</td>
<td>4255</td>
<td>84</td>
<td>3</td>
<td>0.71</td>
<td>2440</td>
<td>1049</td>
<td>906</td>
<td>422</td>
</tr>
</tbody>
</table>

Note: Column 4 is the total number of passed test cases under our ComMT test. We marked the value with underline when the total number of passed test cases is less than the entire test suite size.
While the compression techniques select only a subset of the
mutants for strong mutation, they can also estimate whether the
non-selected mutants are killable by leveraging the groupings
generated by FCA. Therefore, we use the accuracy as further
performance metric, which is defined as follows:

$$\text{accuracy}(C,T) = \frac{(TP + TN)}{\text{total}}$$

(2)

where TP denotes the number of mutants that are strongly
killed by T and that are also correctly identified by a given
method M (true positives); TN is the number of mutants that
are not strongly killed by T and that are correctly identified by
M (true negatives); total denotes the total number of mutants
for the class C. To ease the comparison, we use the mean
and standard deviation of absolute error and accuracy scores
obtained for all classes in a given Java project.

For RQ2 we consider the speed-up metric. When estab-
lishing the speed-up, we should first consider the overhead
induced by an approach. For random sampling, we consider
the overhead to be zero, as mutation sampling does not
require any prerequisite knowledge. For weak mutation, the
overhead consists of one single execution of the test suite
against an instrumented version of the original program. For
the compression strategies, the overhead is composed of the
overhead incurred by both weak mutation and the compression
procedure (mutant clustering and mutant selection):

$$\text{overhead} = \text{exec}\_\text{time}(\text{weak}\_\text{mutation} + \text{compression})$$

(3)

The speed-up metric itself is computed using strong mu-
tation with coverage-based optimisation as the baseline.
We explicitly chose this optimisation as it is already in-
tegrated into several existing mutation testing tools (e.g.,
JAVALANCHE [21], Major [22] and PIT/PITest [23]). The
results of random sampling and compression strategies are the
average values over 100 runs; for weak mutation, the execution
time is zero. Then, the speed-up is defined as follows:

$$\text{speed-up} = \frac{\text{exec}\_\text{time}(\text{strong}\_\text{mutation})}{\text{exec}\_\text{time}(M)}$$

(4)

where the denominator is the execution time of a method M
computed as the sum of its overhead and the execution time
needed to run the tests against the selected mutants.

For RQ3, we first provide a graphical comparison among
the different mutation strategies by using the speedup-error
graphs. In such a graph, the X-axis denotes the speed-up
scores and the Y-axis shows the mean absolute error achieved
by each strategy and for each project in our study; an “ideal”
score would have a high X-value and a low Y-value. We also
use the speedup-accuracy graphs, which plot the speed-up
(mean) on the X-axis and the corresponding accuracy (mean)
on the Y-axis; an “ideal” score would have a high X-value
and a high Y-value. Although we may see some trends via
graphical analysis, we would like to know which strategy
achieves the best speed-up when accepting a given absolute
error rate. Therefore, we consider the following absolute error
thresholds: $\sigma_{e_1} = 5\%$, $\sigma_{e_2} = 10\%$, and $\sigma_{e_3} = 15\%$. Then,
for each threshold $\sigma_{e_i}$ and for each mutation strategy M, we
count the number of projects for which M achieves the highest
speed-up compared to the other strategies while yielding an
absolute error score lower than $\sigma_i$. Similarly, we also consider
three accuracy thresholds, i.e., $\sigma_{a_1} = 95\%$, $\sigma_{a_2} = 90\%$, and
$\sigma_{a_3} = 85\%$. Different from the absolute error, we count
the number of projects in which M achieves an accuracy higher
than the threshold $\sigma_{a_i}$.

**Statistical Analysis.** To assess whether the differences
among the various mutation strategies are statistically signif-
ificant or not, we adopt Friedman’s test [34] with $\alpha = 0.05$.
It is a non-parametric test for comparing multiple treatments
(mutation strategies) in the context of a multiple-problem
analysis (i.e., multiple projects) [34]; it does not require data
to be normally distributed and it is widely applied to compare
randomised algorithms [35], [36] (e.g., random sampling).
While Friedman’s test reveals whether data distributions differ
statistically, tests for pairwise comparison are needed to deter-
mine which treatment outperforms the others. For this, we use
Conover’s post-hoc procedure [37] and we further adjusted the
obtained $p$-values using Holm-Bonferroni [38].

V. RESULTS

A. RQ1: accuracy

Table II reports the mean accuracy and absolute error scores
for each project in our study as well as the corresponding
standard deviation scores.

**Absolute error.** Focusing on the absolute error we observe
that weak mutation performs worst with an error rate of 23% and
a standard deviation of 22% on average. This result is
due to this strategy using all mutants that are weakly covered
to approximate strong mutation coverage, yet, not all infected
states propagate to changes observable in assertions.

Moreover, $\text{overlap+mloc}$ produces the lowest mean absolute
error in 19 out of 20 projects, followed by $\text{fca+mloc}$ and $\text{overlap+mop}$. Overall, the strategies using overlapped grouping
perform slightly better than FCA-based maximal groupings. To
ease the comparison between the compression techniques and
random sampling, in Table II we underline the mean and std
values of a compression technique if they are smaller (better)
than the values of random sampling in the same project. We
can see that the six compression strategies outperform random
sampling in terms of absolute error scores for most projects.
The best (lowest) scores are obtained using the $\text{overlap+mop}$,$\text{overlap+mloc}, \text{fca+mop}$ and $\text{fca+mloc}$ approaches.

Looking at the six compression strategies, we can see that
the mean absolute error and standard deviation scores are
reduced when incorporating knowledge of mutation operators
(mop) and location (mloc). For example, $\text{overlap}$ achieves
an error rate of 15% for the $\text{jpacman}$ project, while when
adding mutation operator and location knowledge it achieves
an error rate of 10% and 7%, respectively. In general, we
find that mutant location is more important than mutation
operator, as the absolute error decreases by 4% on average,
with a minimum improvement of 1% for $\text{assertj}$ and a
maximum improvement of 13% for $\text{zt-}$. This is an
### TABLE I

**SUMMARY OF THE RESULTS FOR RQ1**

<table>
<thead>
<tr>
<th>PID</th>
<th>Overlap</th>
<th>FCA</th>
<th>Random</th>
<th>Weak</th>
<th>PCA</th>
<th>Random</th>
<th>Weak</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Simple</td>
<td>w/ Mop</td>
<td>w/ Mloc</td>
<td></td>
<td>Simple</td>
<td>w/ Mop</td>
<td>w/ Mloc</td>
</tr>
<tr>
<td>1</td>
<td>0.09/1.17</td>
<td>0.08/0.16</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.17</td>
<td>0.11/0.16</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>2</td>
<td>0.07/0.11</td>
<td>0.08/0.11</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>3</td>
<td>0.07/0.12</td>
<td>0.08/0.12</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>4</td>
<td>0.08/0.14</td>
<td>0.08/0.14</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>5</td>
<td>0.09/0.15</td>
<td>0.08/0.15</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>6</td>
<td>0.08/0.16</td>
<td>0.08/0.16</td>
<td>0.08/0.15</td>
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<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>7</td>
<td>0.09/0.17</td>
<td>0.08/0.17</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>8</td>
<td>0.09/0.18</td>
<td>0.08/0.18</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>9</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
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<td>10</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
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<td>11</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>12</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>13</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>14</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>15</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>16</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>17</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>18</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>19</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>20</td>
<td>0.10/0.19</td>
<td>0.08/0.19</td>
<td>0.08/0.15</td>
<td></td>
<td>0.13/0.16</td>
<td>0.11/0.15</td>
<td>0.11/0.16</td>
</tr>
<tr>
<td>Mean</td>
<td>0.09/0.13</td>
<td>0.08/0.15</td>
<td>0.08/0.16</td>
<td>0.13/0.14</td>
<td>0.10/0.16</td>
<td>0.11/0.15</td>
<td>0.13/0.15</td>
</tr>
</tbody>
</table>

### TABLE III

**RANKING PRODUCED BY THE FRIEDMAN’S (LARGER RANK INDICATES SMALLER ERROR) AND STATISTICAL SIGNIFICANCE BY THE CONOVER’S POST-HOC PROCEDURE.**

<table>
<thead>
<tr>
<th>ID</th>
<th>Algorithms</th>
<th>Rank</th>
<th>Significantly better than</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>overlap + Mloc</td>
<td>8.00</td>
<td>(2), (3), (4), (5), (6), (7), (8)</td>
</tr>
<tr>
<td>2</td>
<td>FCA + Mloc</td>
<td>6.55</td>
<td>(4), (5), (6), (7), (8)</td>
</tr>
<tr>
<td>3</td>
<td>overlap + Mop</td>
<td>6.30</td>
<td>(4), (5), (6), (7), (8)</td>
</tr>
<tr>
<td>4</td>
<td>FCA + Mop</td>
<td>4.55</td>
<td>(6), (7), (8)</td>
</tr>
<tr>
<td>5</td>
<td>Overlap</td>
<td>4.40</td>
<td>(6), (7), (8)</td>
</tr>
<tr>
<td>6</td>
<td>Random Samp.</td>
<td>2.70</td>
<td>(8)</td>
</tr>
<tr>
<td>7</td>
<td>FCA</td>
<td>2.50</td>
<td>(8)</td>
</tr>
<tr>
<td>8</td>
<td>Weak Mut.</td>
<td>1.10</td>
<td></td>
</tr>
</tbody>
</table>

### TABLE IV

**OVERLAP WITH MUTATION LOCATION KNOWLEDGE OUTPERFORMS ALL OTHER MUTATION STRATEGIES IN TERMS OF BOTH ABSOLUTE ERROR AND PREDICTION ACCURACY.** Random sampling is statistically worse than all mutant compression techniques.

**B. RQ2: speed-up**

**Speed-up performance.** Table IV summarises the overall speed-up for the eight approaches in our comparison. For each project, we highlight the two strategies achieving the best speed-up scores in **bold**. Notice that speed-up measures the overall execution time of strong mutation divided by the overall execution time of a mutation strategy. Hence, higher values denote a larger improvement in execution time.
We observe that weak mutation shows the highest speed-up scores since it requires only one test suite execution compared to the original program. Except for weak mutation, FCA achieves the highest speed-up scores in 19 out of 20 cases. FCA is also faster than random sampling, which selects 10% mutants for strong mutation. Indeed, the former is 6.6 times faster than the latter on average, with a minimum speed-up of 2.6X (in vраптор) and a maximum one of 15.4X (in stream-1lb). This is because FCA suggests on average less than 10% of mutants (with a minimum of 0.4% of mutants) to evaluate in strong mutation analysis. Instead, the sampling strategy constantly (and randomly) selects 10% of mutants to execute.

The only exception to the previous finding is represented by distributedlog for which random sampling is faster than FCA. In this case, the total percentage of mutants that are injected into statements covered by the test suite (reachability condition) is fairly low, being 2.8%. Thus, random sampling can achieve a considerable speed-up if we leverage the coverage-based optimisation, i.e., if we skip uncovered mutants (i.e., mutants of uncovered statements). Instead, FCA selects almost twice as many mutants for these projects.

To further ease the comparison, in Table IV we underline the compression strategies that achieve better speed-up scores than random sampling in each project. We observe that overlap, FCA+mop and overlap+mop outperform random sampling in terms of speed-up in 19 projects out of 20. On average, they are respectively 3.7X, 3.0X, and 2.1X faster than random sampling. It is worth noticing that the number of selected mutants does not directly determine the overall speed-up. For example, for the project pac4j, overlap selected 21.0% of mutants, which is larger than the percentage of mutants selected by mutation sampling (i.e., 10%). However, overlap achieves a larger speed-up of 4.8X against 3.1X of mutation sampling. The reason is that compression strategies use weak mutation information to further filter out the unnecessary test executions, while mutation sampling does not.

From the comparison of the six compression strategies, we observe that including mutant location leads to selecting more mutants for strong mutation, thus, reducing the overall speed-up. For example, overlap+mloc achieves lower speed-up scores than overlap in all 20 projects. Moreover, by comparing the two strategies based on mutation location knowledge (i.e., overlap+mloc, and fca+mloc) with random sampling, we observe that the differences in terms of speed-up are small. Indeed, the average speed-up scores of overlap+mloc, fca+mloc and random sampling are 5.13, 5.44 and 5.42, respectively. Instead, selecting mutants according to mutation operator generates a lower number of mutants to evaluate in strong mutation compared to mutation location.

Our findings are confirmed by Friedman’s test: the mutation strategies are statistically different in terms of speed-up scores (p-value = 10^-16). According to Conover’s procedure, weak mutation and FCA statistically outperform all other mutation strategies. Moreover, random sampling is ranked sixth and is statistically more efficient than overlap+mloc only, although the difference is marginal as suggested by the average scores reported in Table IV. Instead, FCA, overlap, FCA+mop and overlap+mop are statistically superior to random sampling.

### Overhead
In the previous paragraphs, we observed that most of mutation compression strategies are more efficient than random sampling. Here, we investigate the overhead that is due to the different steps that such strategies implement. In the following, we consider as compression overhead the overlap and random sampling.
execution time needed to compute the maximal and/or overlapped groupings; while the overall overhead is the sum of the compression overhead, the time for running all tests once for weak mutation, and the time to select the mutants.

Table IV reports the compression overhead and overall overhead as a percentage (ratio) of the full execution time of each strategy, which also includes the time needed to run the selected tests and mutants for strong mutation. The highest values for each project are highlighted in bold face. From Table IV, we can observe that the compression overhead takes up less than 0.001% of the total execution time; thus, it is negligible with respect to the execution time of evaluating the selected mutants for strong mutation. The overall overhead accounts for up to 9.30% of the total execution time and weak mutation represents the larger portion of this overhead. Among the 20 projects, the overall overhead of the strategy FCA is likely higher than the other compression strategies. However, the differences among them are lower than 0.1%.

Weak mutation scores best among the eight techniques in terms of speed-up. Without considering weak mutation, four mutant compression strategies are statistically more efficient (have better speed-up scores) then random sampling.

C. RQ3: trade-offs

From the results of RQ1 and RQ2, it is clear that the mutation strategies that perform best in terms of accuracy are also the more expensive to perform. Weak mutation and FCA grouping strategies perform best in terms of speed-up, while the overlapped grouping strategies, and overlap+mloc in particular, better approximate the strong mutation score. Therefore, in this section, we analyse the trade-offs between speed-up, absolute error and accuracy.

From Figure 3, we observe that weak mutation achieves the best speed-up, while its absolute error and accuracy typically score worst when compared to other techniques. FCA comes second in terms of the overall speed-up, but its absolute error as well as accuracy are better than weak mutation for most of the projects. We notice that overlap is slightly slower than FCA, but shows a small improvement in both absolute error and accuracy when compared to FCA. FCA+mop and overlap+mop have quite similar trade-offs considering speed-up and absolute error as their data points are very close to each other. However, in terms of speed-up and accuracy, fca+mop is slightly faster than overlap+mop, while overlap+mop is more accurate than fca+mop. Moreover, overlap+mloc, fca+mloc and random sampling have the same speed-up score; however, the absolute error of random sampling is higher than that for the other two strategies. Overlap+mloc and fca+mloc are the most accurate strategies in terms of both absolute error and accuracy, but their speed-up performance is the least good.

Table VI shows the number of projects for which each strategy Mi provides the best speed-up score at different thresholds.

![Graphical comparison of the eight mutation strategies in terms of speed-up, absolute error and accuracy.](image)

**TABLE VI**

<table>
<thead>
<tr>
<th>Absolute Error</th>
<th>Overlap</th>
<th>FCA</th>
<th>Random</th>
<th>Weak</th>
</tr>
</thead>
<tbody>
<tr>
<td>Simple</td>
<td>w/ Mop</td>
<td>w/ Mloc</td>
<td>Simple</td>
<td>w/ Mop</td>
</tr>
<tr>
<td>≤5%</td>
<td>2</td>
<td>4</td>
<td>10</td>
<td>0</td>
</tr>
<tr>
<td>≤10%</td>
<td>4</td>
<td>1</td>
<td>1</td>
<td>5</td>
</tr>
<tr>
<td>≤15%</td>
<td>1</td>
<td>0</td>
<td>1</td>
<td>11</td>
</tr>
<tr>
<td>overall</td>
<td>7</td>
<td>5</td>
<td>12</td>
<td>16</td>
</tr>
</tbody>
</table>

![Graphical representation of the trade-offs between speed-up, absolute error and accuracy.](image)

(a) error rate

<table>
<thead>
<tr>
<th>Accuracy</th>
<th>Overlap</th>
<th>FCA</th>
<th>Random</th>
<th>Weak</th>
</tr>
</thead>
<tbody>
<tr>
<td>Simple</td>
<td>w/ Mop</td>
<td>w/ Mloc</td>
<td>Simple</td>
<td>w/ Mop</td>
</tr>
<tr>
<td>≥95%</td>
<td>2</td>
<td>2</td>
<td>9</td>
<td>0</td>
</tr>
<tr>
<td>≥90%</td>
<td>5</td>
<td>4</td>
<td>2</td>
<td>3</td>
</tr>
<tr>
<td>≥85%</td>
<td>4</td>
<td>2</td>
<td>1</td>
<td>8</td>
</tr>
<tr>
<td>overall</td>
<td>11</td>
<td>8</td>
<td>12</td>
<td>11</td>
</tr>
</tbody>
</table>

(b) accuracy

of 20. This indicates that using mutation location leads to more accurate estimations of the actual strong mutation. When considering a 10% error rate threshold, FCA+mop and FCA show the largest speed-up for the majority of projects. Instead, when the goal is to reach an absolute error ≤15%, FCA has the best speed-up scores for 11 out of 20 projects. Moreover, we notice that random sampling performs worst as it has speed-up scores that are always lower than the other mutation strategies.
at the same (or higher) level of absolute error.

Similar results can be observed when considering different thresholds for accuracy. As reported in Table VI, overlap+mloc and fca+mloc show the best speed-up for accuracy ≥95% for the relative majority of the projects. When considering an accuracy ≥90%, overlap and fca+mloc show the best speed-up for 5 out of 20 projects each. Instead, if we focus on 85% of accuracy, we observe that FCA is the best approach to choose as it shows a speed-up ranging from 6.3X up to 94.3X compared to strong mutation. Finally, if we focus on 85% accuracy, we observe that FCA is the best approach to choose.

Overlap+mloc provides the best speed-up scores when the goal is to achieve an accuracy >95% or an absolute error <0.05. Other mutation compression strategies provide larger speed-up, but with a corresponding decrease in accuracy. Random sampling is less accurate and/or slower than all compression strategies.

D. Discussion

Looking at all the results, we can observe that random sampling with 10% sampling ratio is able to speed up (strong) mutation testing from 2.0 to 22.7X with an absolute error within 15% for 80% of the projects. Mutation sampling is also easy to apply in mutation tools as it does not require any prerequisite knowledge of the program context and mutation operators. However, mutation strategies based on compression techniques achieve better speed-up (i.e., are more efficient) and/or lower absolute error than random sampling. For example, overlap+mloc yields an absolute error which is always lower than 9% with a speed-up ranging between 2.0 to 53X. This represents an important finding if we consider the recent study by Gopinath et al. [17], which showed that other mutation reduction techniques (including e-selective) provide small or negligible improvements in effectiveness and are more expensive compared to random sampling.

Another disadvantage of random sampling is that it can estimate the overall mutation score, but cannot estimate whether each mutant is strongly killable or not (it only does for the sampled mutants). The overall mutation score is of course very important when assessing the test suite quality at a high level; however, Coles [39, slide 57] observed that programmers prefer to obtain specific insights into which mutants their test suite is able to kill. From this perspective, mutation compression strategies select a subset of “representative” mutants for the programmers to investigate. In addition, overlap+mloc and FCA+mloc guarantee that for every possible statement that can be mutated, at least one mutant will be selected. Although this may negatively affect the speed-up compared to random sampling, programmers can benefit from the killable mutant results at every possible mutant location.

VI. Threats to Validity

Threats to external validity: Our results are based on mutants generated by the operators implemented in EvoSuite; these results might be different when using other mutation tools [40]. With regard to the subject selection, we chose 18 out of the 20 projects from GitHub’s top starred 3000 repositories; the selected projects differ in size, number of test cases and application domain.

Threats to internal validity: The main threat for our study is the implementation of the compression strategies. For FCA, we use its implementation available in MATLAB [41], which is a well-known scientific software. For the instrumentation and the mutation operators, we relied on their implementation available in EvoSuite [42]. Moreover, we carefully reviewed and tested all code for our study to eliminate potential faults in our implementation.

Threats to construct validity: The main threat is the measurement we used to evaluate our methods. We minimise this risk by adopting evaluation metrics that are widely used in research, as well as proper statistical analysis to assess the significance.

VII. Conclusions

In this paper, we have conducted a detailed investigation of different compression techniques to speed up mutation testing based on the work by our previous work [18]. We have enhanced our original FCA-based compression strategy in two distinct ways: (1) by proposing a novel mutant clustering algorithm, overlapped grouping, in addition to FCA; (2) by incorporating mutation location and mutation operator information in the compression procedure. Thereby, we have introduced and investigated six compression strategies based on two clustering algorithms and three mutant selection strategies.

The results of an empirical study with 20 open-source projects show that mutant compression techniques can effectively speed up strong mutation testing up to 94.3 times with an accuracy >90%. FCA is the fastest strategy while overlap + mloc is the most accurate. In comparison, weak mutation attains a higher absolute error (23%) and lower accuracy (75%). Random sampling with 10% as sampling percentage is statistically less accurate than all mutant compression strategies, and worse in terms of speed-up than four compression strategies (excluding the two with knowledge of mutation locations).

Another important finding is that mutation location trumps mutation operator information when selecting mutants to evaluate for strong mutation. Hence, researchers should take into account the mutation location in addition to the mutation operators when detecting redundant or subsuming mutants (e.g. [25]–[27]). This is a clear invitation for future work.

Since our results are encouraging, we envision the following future work: (i) combining mutant location and mutation operator information; (ii) investigating other compression methods, such as Principal Component Analysis [43]; (iii) applying compression techniques in mutation-based test case generation [44], [45].
REFERENCES


