A Moment of Weakness: Protecting Against Targeted Attacks Following a Natural Disaster

Jorik Oostenbrink  
Delft University of Technology  
J.Oostenbrink@tudelft.nl

Fernando A. Kuipers  
Delft University of Technology  
F.A.Kuipers@tudelft.nl

ABSTRACT

By targeting communication and power networks, malicious actors can significantly disrupt our society. As networks are more vulnerable after a natural disaster, this moment of weakness may be exploited to disrupt the network even further. However, the potential impact and mitigation of such a follow-up attack has yet to be studied.

In this paper, we propose a framework to analyze the impact of a combination of a natural disaster followed by a targeted single node failure. We apply this framework on empirical disaster data and two network topologies. Our experiments show that even small targeted attacks can significantly augment the already grave network disruption caused by a natural disaster. We further show that this effect can be mitigated by adopting a calculated repair strategy.

1. INTRODUCTION

Communication and power networks are critical to our society. This makes them a prime target for malicious actors trying to destabilize or terrorize a country. In fact, Admiral Mike Rogers, the former director of the NSA, warned, “It’s only a matter of the when, not the if, you are going to see a nation state, a group or an actor engage in destructive behavior against critical infrastructure of the United States” [1].

For many of these actors, the exact timing of their attack may not be essential. Their focus is to deal as much damage as possible, preferably using only a small amount of resources. A strategy they might adopt is to delay their attack until the network is most vulnerable, such as after a natural disaster. By attacking the network at its weakest moment, a bad actor can multiply the damage he, or the disaster by itself, could otherwise inflict.

As it takes time to prepare and execute an attack, a network operator has a limited window to try to reduce the impact of any potential attack. However, to the best of our knowledge, while there is a large body of research on the resilience of networks to natural disasters (see [4]) and targeted attacks (e.g. [2, 5, 8]), the potential combination of a disaster followed by a targeted attack has yet to be studied.

In this paper, we propose a framework to analyze the impact of a combination of a natural disaster and targeted single node failure. Our main contributions are as follows:

- We extend our successive disaster framework from [7] to incorporate targeted attacks.
- We apply our framework to empirical disaster data and show that a small follow-up attack can significantly increase the impact of a natural disaster. In addition, we study the effect of changing the repair strategy to prepare for potential follow-up attacks.

2. FRAMEWORK

In [7], we have introduced a model and framework for assessing the resilience of networks to successive disasters, taking into account network recovery. In this paper, we extend this framework to include the risk of targeted attacks.

We model the network as a directed multigraph $G = (V, E, \psi)$ with nodes $v \in V$ connected by links $e \in E$, where $\psi: E \rightarrow V \times V$ and $e \in E$ connects $v_1$ to $v_2$ if and only if $\psi(e) = (v_1, v_2)$. We define the network state $s$ of this network by its failures: network component $c \in V \cup E$ is functioning if and only if $c \notin s$.

Now, to be able to assess the resilience of this network to follow-up attacks, we first need to model the impact of a disaster itself. We assume disaster occurrences are Poissonian, and we are given a multiset $D^*$ of disaster processes $d = (a_d, \lambda_d)$, where $a_d \subseteq V \cup E$ are the components affected by $d$ and $\lambda_d$ is the rate of $d$. Thus, if disaster $d \in D^*$ occurs at time $t$, when the network state is $s$, the new network state at time $t$ will be $s \cup a_d$. As the combination of multiple Poisson processes is itself Poissonian, the disaster processes in $D^*$ can be combined as follows:

$$D = \{(a_d, \lambda_d) | a_d \neq \emptyset \land \lambda_d = \sum_{(a_d, \lambda_d^*) \in D^*} \lambda_{d^*} > 0\}$$ (1)

In this paper, we only consider attacks after a single disaster. In other words, we assume a single disaster occurs and is then followed by an attack on the network. However, by applying the techniques of [7], our model can be easily generalized to capture an attack after an arbitrary number of successive disasters or any other mix of natural disasters and attacks. We fix the time of the initial disaster ($D_1$) at $T_1 = 0$. Now, we can compute the distribution of the network state at $T_1$, $S_1$, by

$$P(S_1 = s) = \sum_{d \in D | a_d = s} P(D_1 = d) = \sum_{d \in D | a_d = s} \lambda_d \lambda_D$$ (2)

where $\lambda_D$ is $\sum_{d \in D} \lambda_d$.

A follow-up attack after a disaster can be pre-planned or opportunistic. In either case, it will take some time to react.
to the disaster and execute the attack. We consider two different attack models: (1) the attack occurs after a fixed amount of time \( t_{\text{attack}} \), and (2) the time between the disaster and attack is exponentially distributed with rate \( \lambda_{\text{attack}} \). In both cases, if the network has been fully repaired before the attack has been executed, we assume the attack will be canceled and the network will not suffer any further damage.

Let \( T_{\text{attack}} \) be the time of the attack. We assume the attacker has perfect knowledge of the network at all times, and will always take down the node that maximizes the number of disconnected node pairs at \( T_{\text{attack}} \). In other words, an attack is modeled as a worst-case node failure.

The target and impact of this attack greatly depend on the progress of network repair at \( T_{\text{attack}} \). We consider a deterministic repair model. That is, we assume that, given a certain starting state, the recovery of the network is fixed (until the attack occurs). For each possible starting state, we define a repair function \( r_s \). For any further damage.

\[
\lambda_d \text{ failed nodes} \rightarrow \lambda_{\text{attack}} \text{ attack} \rightarrow V \cup E.
\]

For Sinet, we consider both types of disasters, while for Deltacom, we only consider tropical cyclones. We assume only network nodes are affected by these disasters and all network links remain functioning. This gives us a yearly disaster rate \( \lambda_d \) of 1.597 for Sinet and 1.342 for Deltacom.

For ease of reading, we make the assumption that one node is repaired every day. However, by scaling both the attack and repair time, our results can easily be transformed to any other repair time.

We use the Average Two-Terminal Reliability (ATTR) as an impact measure. The ATTR is defined as the number of connected node pairs divided by the total number of node pairs in the network. That is, \( \text{ATTR} = \frac{M}{|V|} \).

3. EXPERIMENTS

In this section, we apply our framework to two slightly modified versions of undirected networks from the topology zoo [6]: Sinet and Deltacom. Sinet is a Japanese network of 47 nodes connected by 49 links, and Deltacom is a US network of 99 nodes connected by 151 links.

We make use of the same disaster set as was used in [7]: a set of earthquake scenarios and historical tropical cyclones. For Sinet, we consider both types of disasters, while for Deltacom, we only consider tropical cyclones. We assume only network nodes are affected by these disasters and all network links remain functioning. This gives us a yearly disaster rate \( \lambda_d \) of 1.597 for Sinet and 1.342 for Deltacom.

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3.1 Impact

We first compare the impact of follow-up attacks to those of a disaster or attack by itself. We assume both networks use a greedy repair function that continuously chooses the node with the largest impact on ATTR to repair. To make a fair comparison, we modify the attack function \( att \) by continuing the follow-up attack even if the network has been fully repaired.

Figure 1 shows the expected ATTR after a targeted attack, disaster, or disaster and follow-up attack. Sinet is clearly more vulnerable to both targeted attacks and disasters than Deltacom. However, for both networks, a follow-up attack can significantly increase the impact of a disaster. For Sinet, the combination of disaster and follow-up attack disconnects more than half of all node pairs on average.

We have assumed that a malicious actor strikes after the first natural disaster that hits the network. On average, such an opportunity occurs more than once per year. However, he could also decide to wait for a larger disaster, which would allow him to inflict even more damage to the network.

We consider an attacker that waits for a disaster that damages at least \( k \) (10) of Deltacom’s (Sinet’s) nodes. Figure 2 shows the impact of this more patient follow-up attack. Waiting for larger events allows the attacker to inflict much more damage. In the case of Sinet, the expected impact of the follow-up attack is the disconnection of around half of the remaining node pairs.

3.2 Repair Strategies

After a disaster, the network operator will typically try to restore as much functionality as quickly as possible. The impact of the follow-up attack greatly depends on the progress of these repair operations at the time of the attack. Although speeding up repair would have the largest effect, the network operator can also change the order in which components are repaired to try to minimize the impact of any

\[ M_{ad,s} \text{ is the first time } t \text{ at which } att(r_{ad}(t)) = s \text{ (or } \infty \text{ if no such time exists)}, \text{ and } S_{ad,s} \text{ is the first time } t \text{ at which } M_{ad,s} \text{ at which } att(r_{ad}(t)) \neq s \text{ (or } \infty \).
In this section, we consider the effect of changing the node repair order on the impact of the follow-up attack. We compare 5 different repair strategies:

- **Betweenness Repair**: Repair nodes in the order of their betweenness centrality [3].
- **Greedy Repair**: Every day, repair the node with the highest impact on the ATTR.
- **Optimal Unprepared Repair Strategy**: Maximizes the ATTR at the time of attack.
- **Greedy Prepared Repair**: Every day, repair the node that would increase the ATTR the most if the network would be attacked immediately afterwards.
Critical infrastructure networks are prime targets for malicious actors trying to destabilize or terrorize a country. As part of their attack strategy, they might wait for points in time when critical infrastructure is significantly more vulnerable, for example right after a natural disaster has struck. However, current disaster vulnerability frameworks do not consider the potential risk of a network to these potential follow-up attacks.

We have proposed a framework for assessing the impact of follow-up attacks. Our framework can take into account a variety of natural disasters and two kinds of attacks: a worst-case node failure after (1) a fixed amount of time or (2) exponentially distributed random delay after an initial disaster.

In our experiments, we have shown that small targeted attacks can significantly augment the impact that a natural disaster has on the network. Fortunately, our results also reveal that the right choice of repair strategy allows network operators to reduce the threat of follow-up attacks at almost no cost to network performance compared to other repair strategies. Our framework aids in determining the efficacy of repair strategies.

5. REFERENCES


